Embark: Securely Outsourcing Middleboxes to the Cloud

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Background

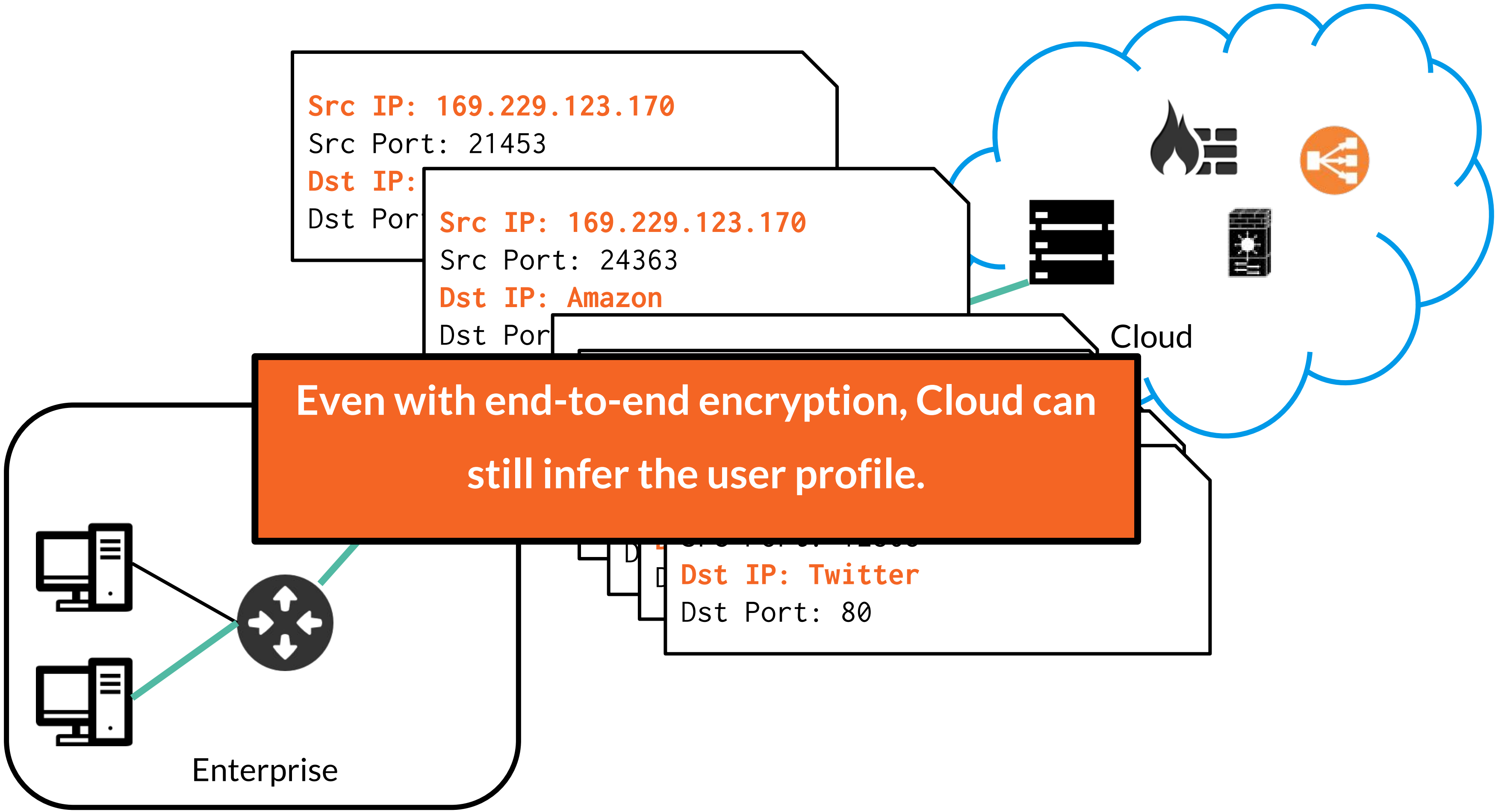
- **Middleboxes are prevalent and problematic**
  - Number of Middleboxes ≈ Number of Routers (APLOMB [SIGCOMM ‘12])
  - Lots of Problems:
    - MB Manifesto [HotNets ‘11], CoMb [NSDI ‘12],
      Honda et al. [IMC'11], DOA [OSDI '04], ETTM [NSDI '11], …

- **A Promising Solution: Outsourcing**
  - APLOMB [SIGCOMM ‘12]
  - Aryaka, Zscaler
  - AT&T NFV/CORD

New Challenge: Confidentiality and Privacy

- **The middleboxes sees the traffic unencrypted.**

- **Strawman: End-to-end Encryption** (e.g. TLS):
  - Some middleboxes cannot process traffic (e.g. Deep Packet Inspection).
  - Unencrypted packet fields still leak information


Src IP: 169.229.123.170
Src Port: 21453
Dst IP:
Src IP: 169.229.123.170
Src Port: 24363
Dst IP: Amazon
Cloud
Even with end-to-end encryption, Cloud can still infer the user profile.
Dst IP: Twitter
Dst Port: 80
Enterprise

Problem Statement

Can we outsource middleboxes without compromising privacy?

Embark

the first system that allows middlebox outsourcing,

while keeping traffic confidential.

Overview

- **Approach**
  - Middleboxes process encrypted traffic without decrypting it
- **Crypto Primitives**
  - **KeywordMatch**: For Signature Matching
    - BlindBox [SIGCOMM ‘15]: Prohibitive Setup Time Per Flow

      Contribution: System Design + Implementation without Per-flow Setup Time
  - **PrefixMatch**: Prefix/Range Matching

    Contribution: A fast, secure encryption scheme for prefix matching

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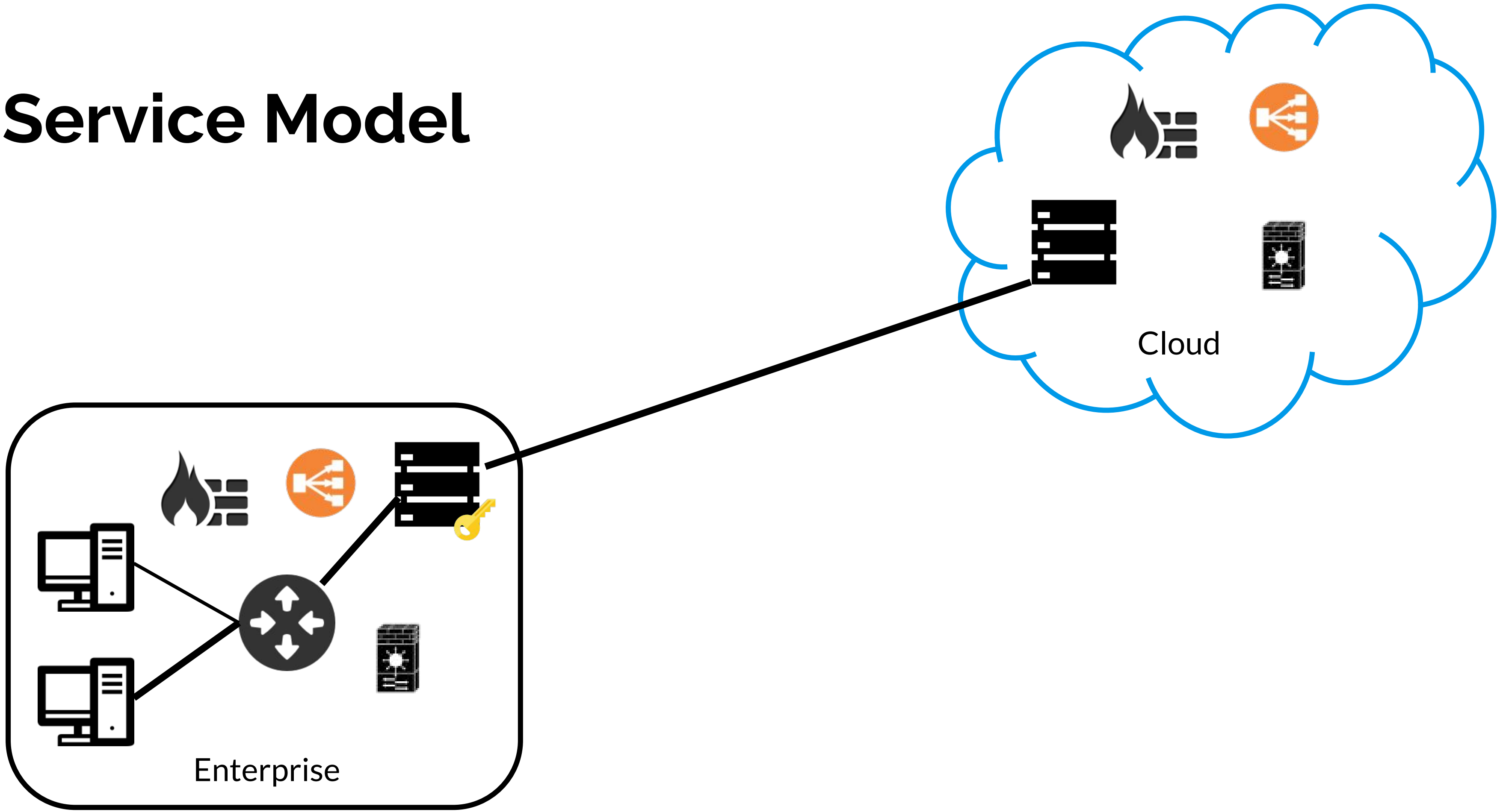
      Contribution: System Design + Implementation without Per-flow Setup Time
  - **PrefixMatch**: Prefix/Range Matching

    Contribution: A fast, secure encryption scheme for prefix matching

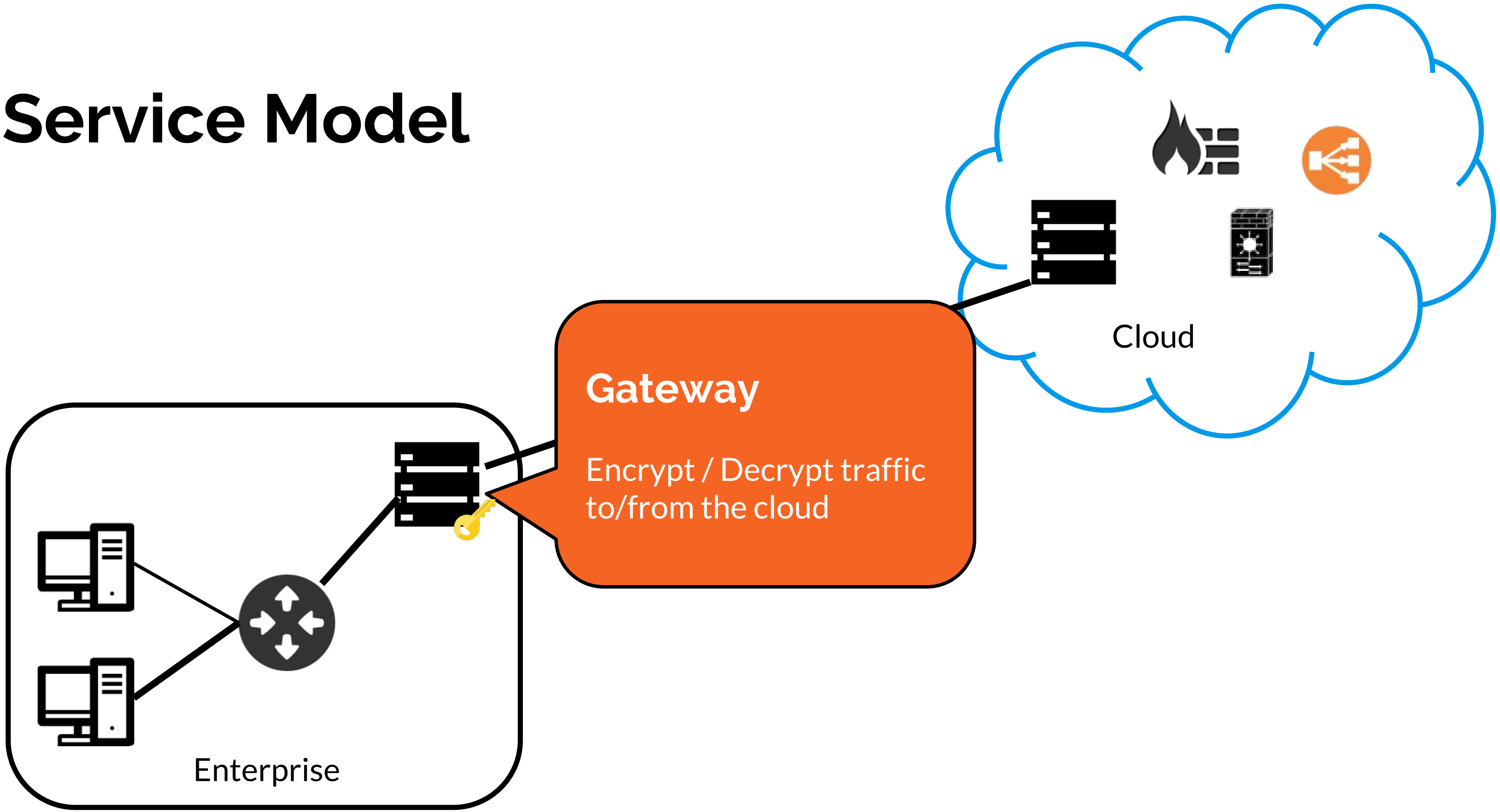
Outline

1. **Service Model of Embark**
2. PrefixMatch: Two Functions
   - **EncryptRanges**
   - **EncryptValue**
3. Evaluation
4. Conclusion

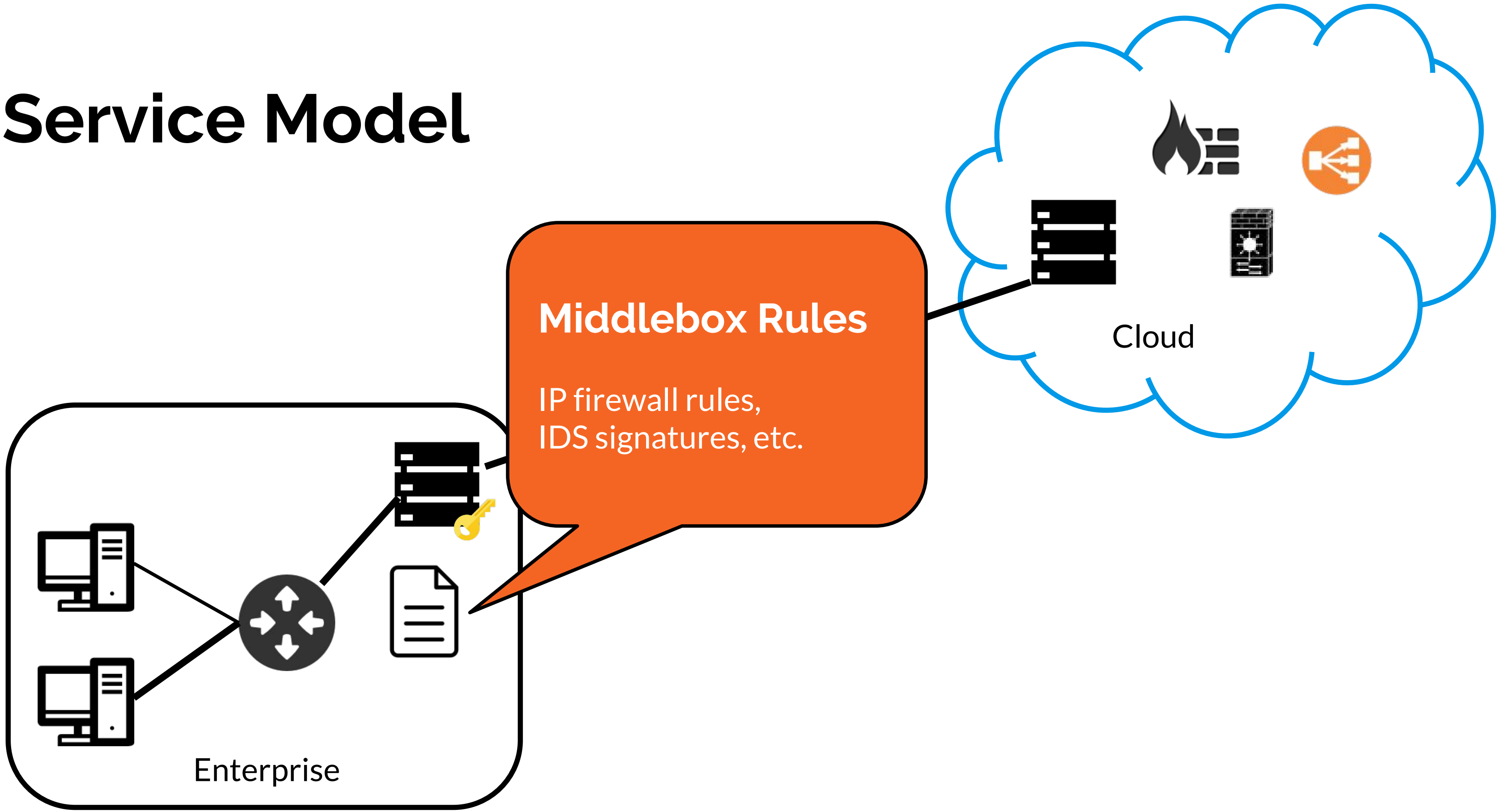
Service Model



Service Model

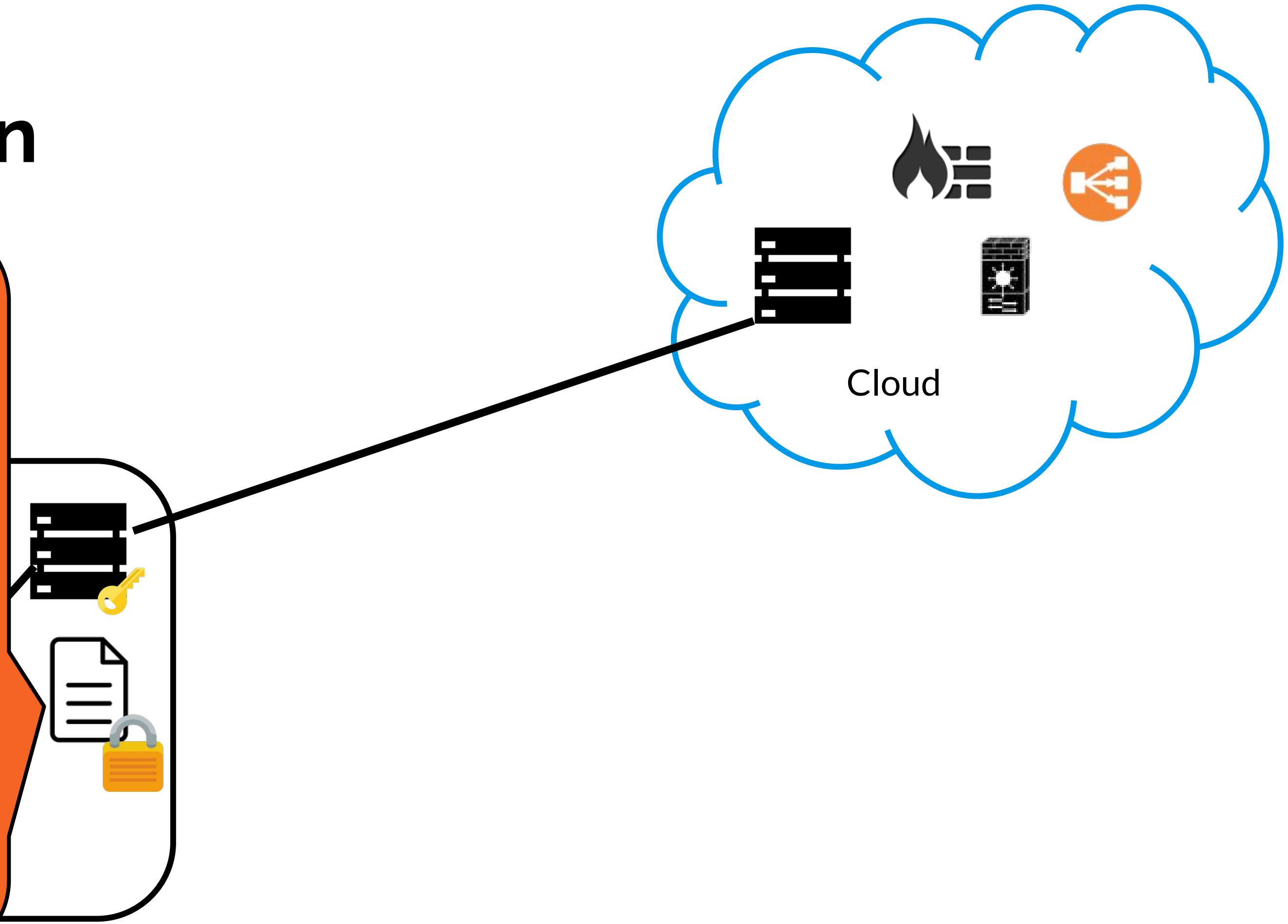


Service Model

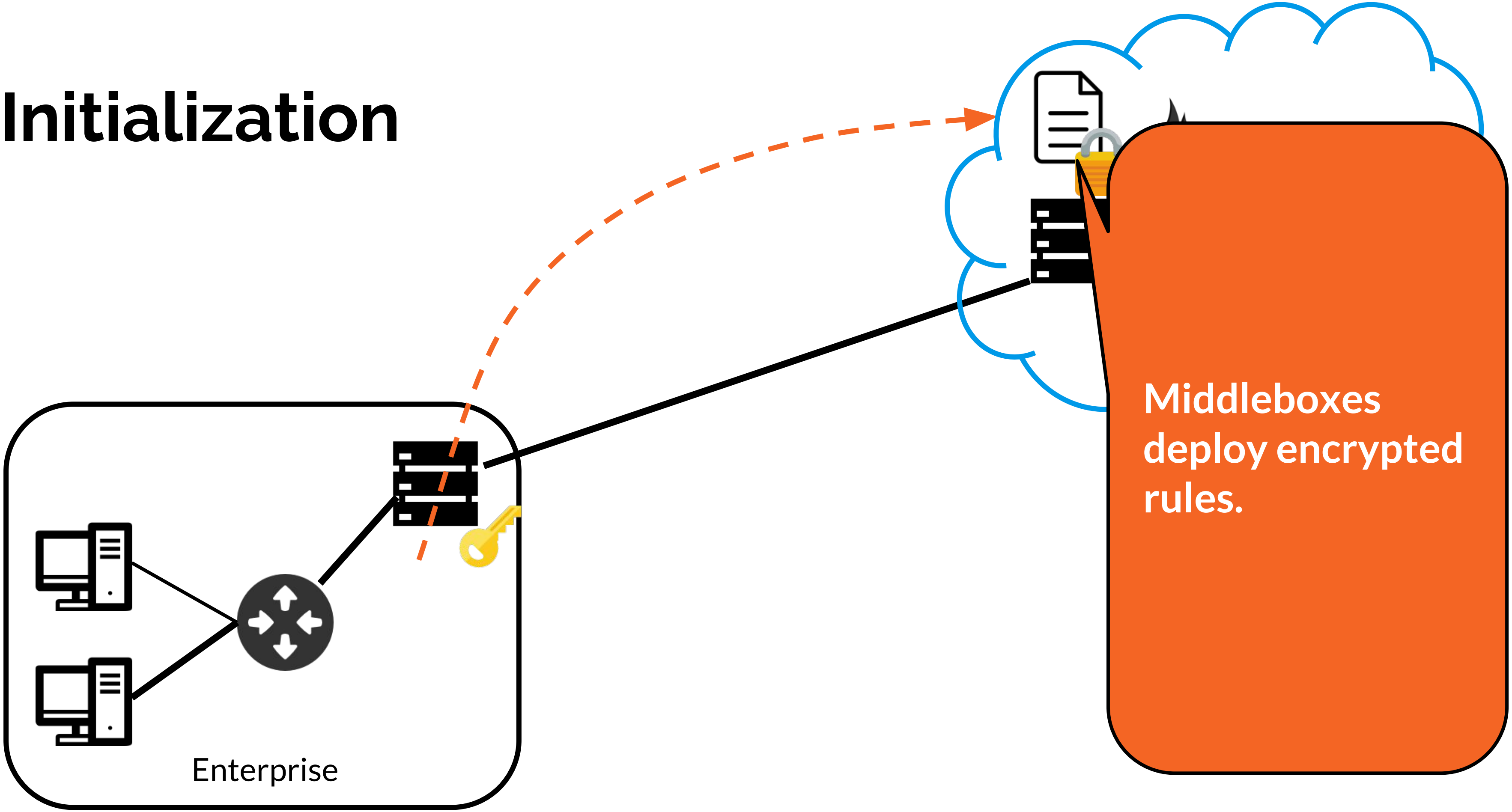


Initialization

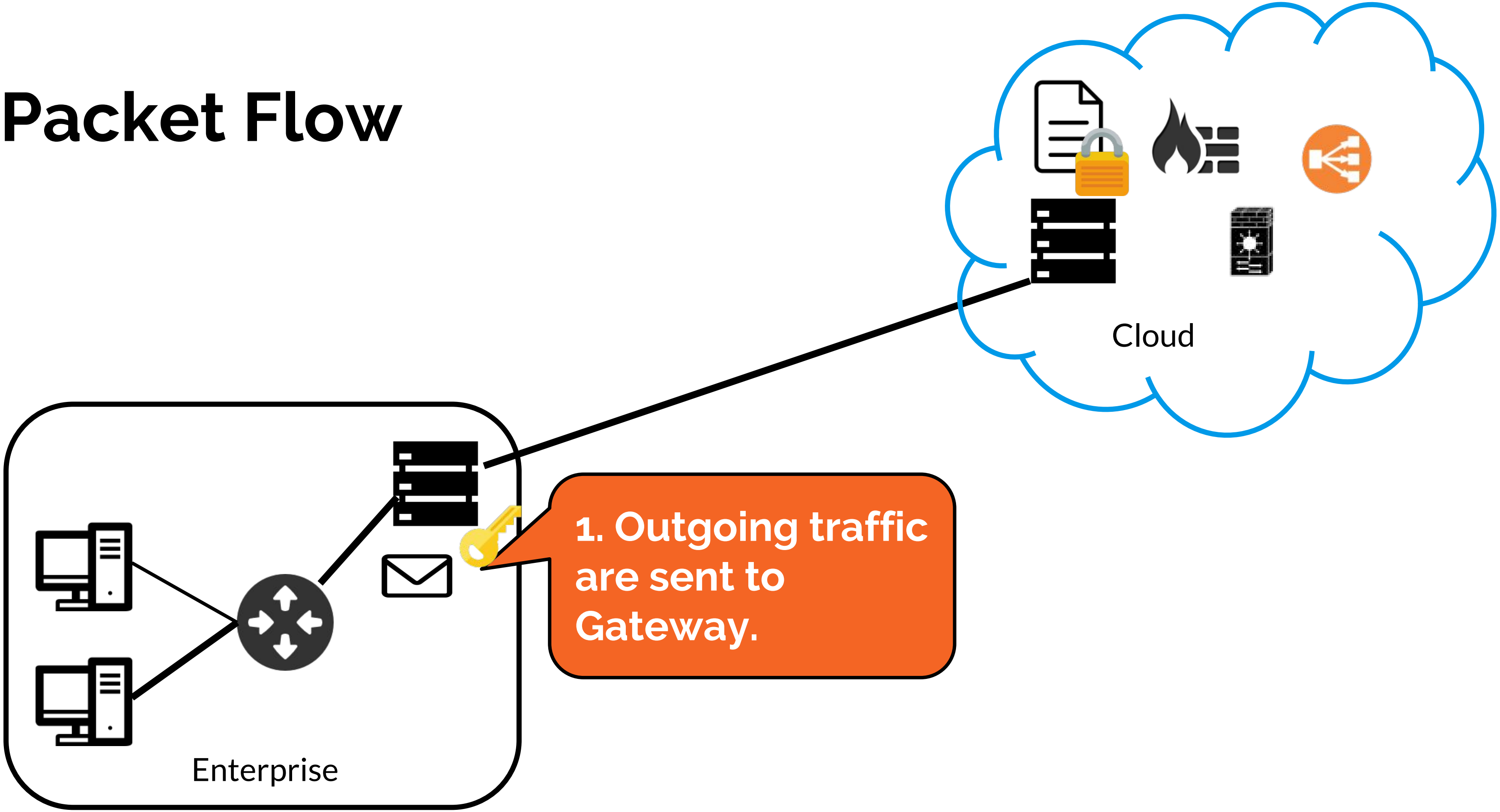
Enterprise encrypt rules using ***EncryptRanges***.



Initialization



Packet Flow

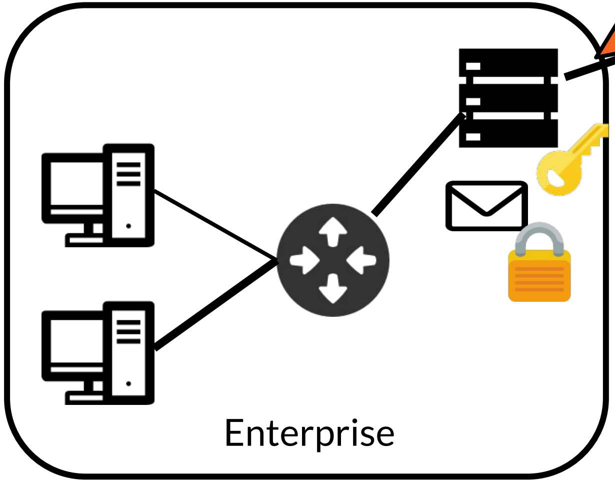


Packet Flow

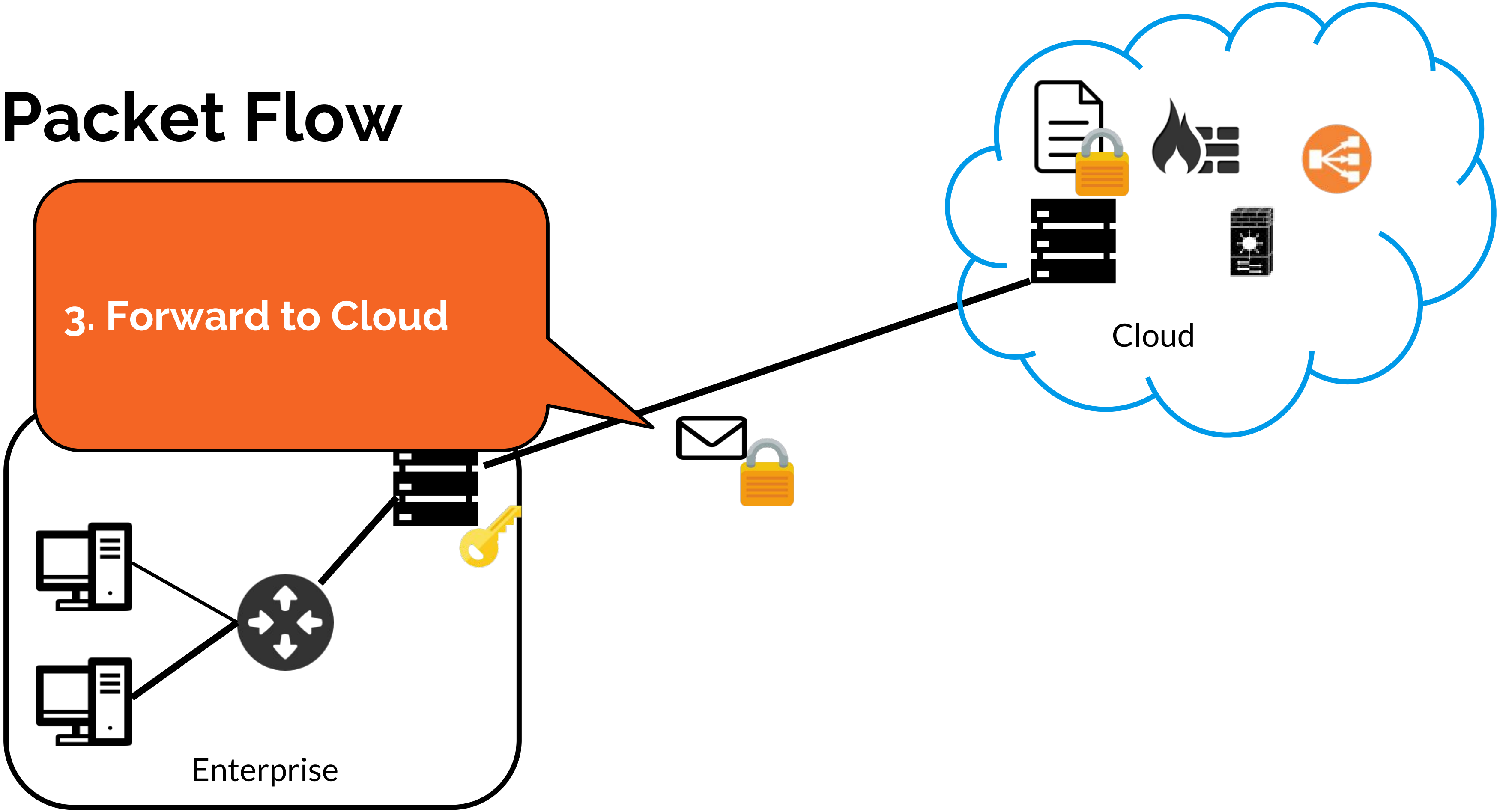
2. Encrypt the traffic

- Encrypt **packet headers** *field by field* using ***EncryptValue***
- Encrypt **payloads** using stream cipher

Implication: no change to packet structure



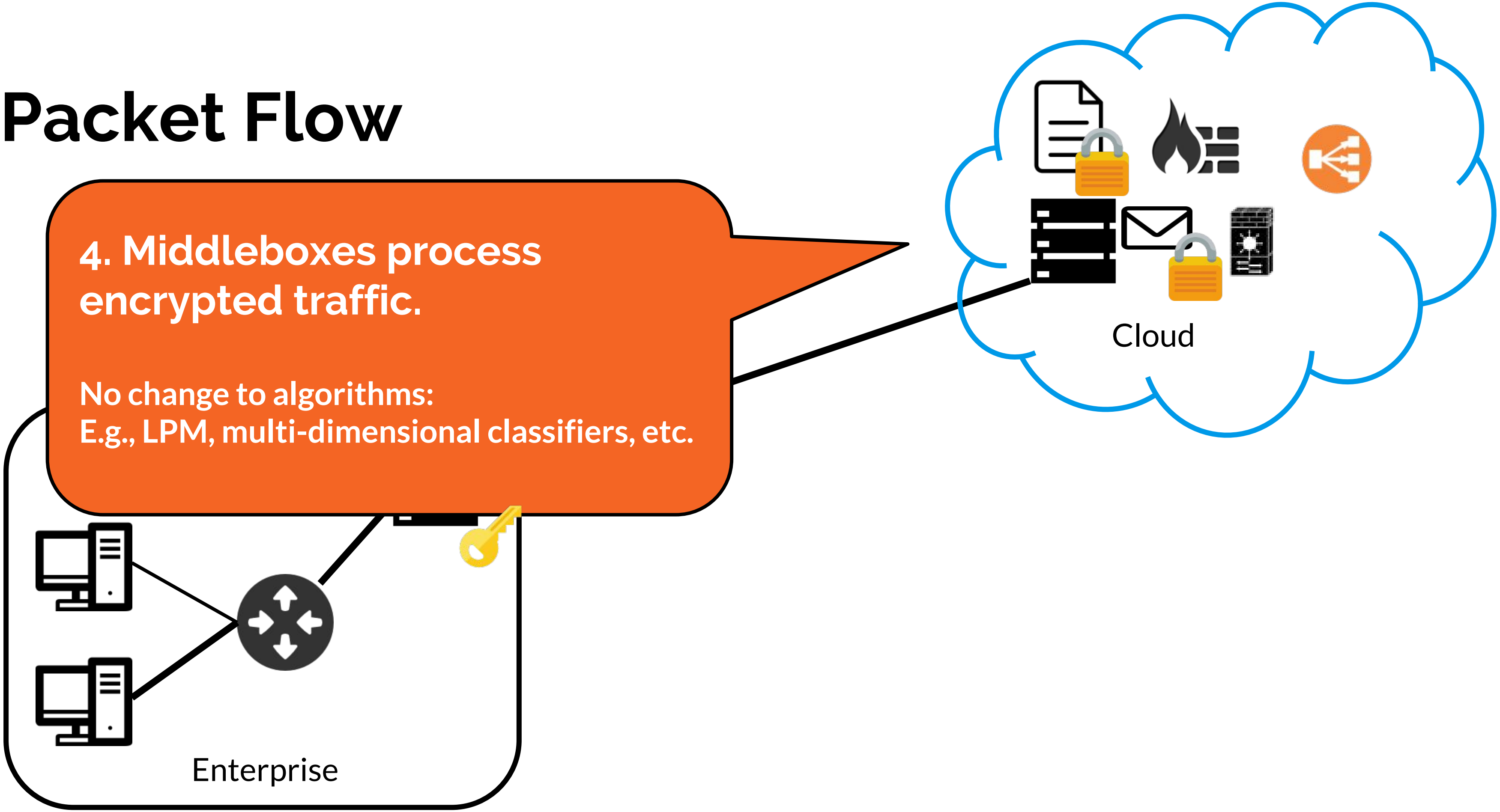
Packet Flow



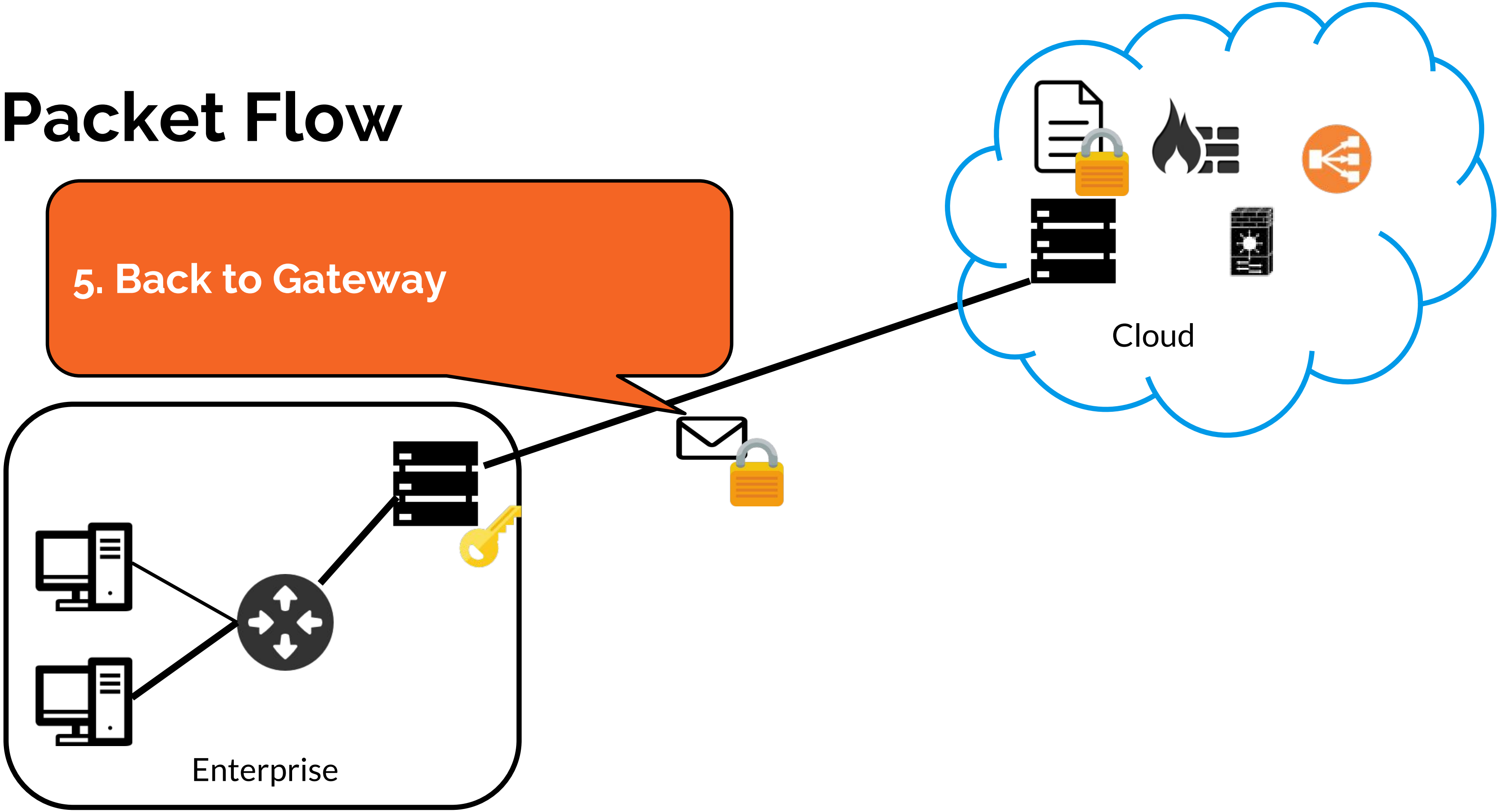
Packet Flow

4. Middleboxes process encrypted traffic.

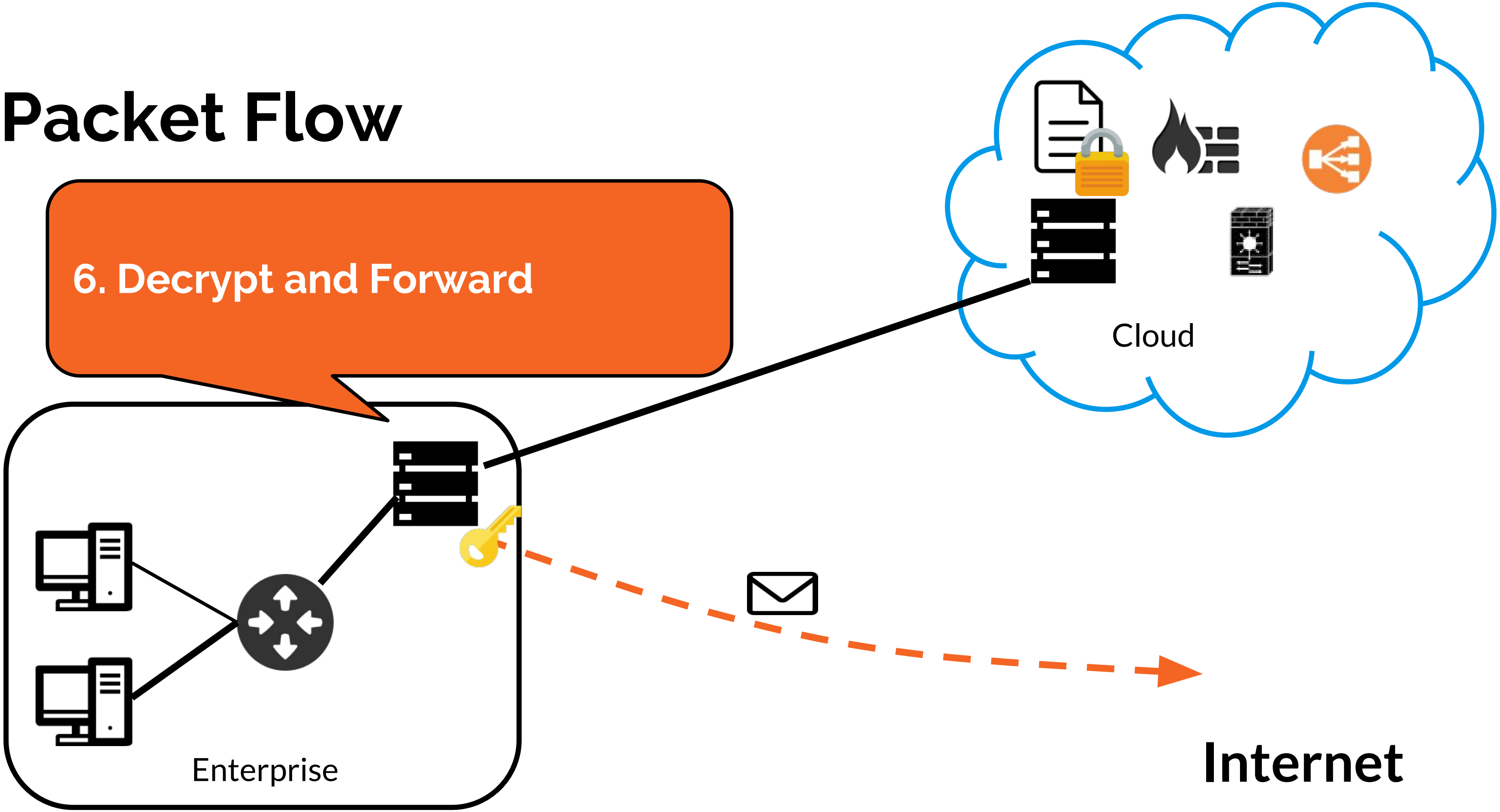
No change to algorithms:
E.g., LPM, multi-dimensional classifiers, etc.



Packet Flow



Packet Flow



Outline

1. Service Model of Embark
2. **PrefixMatch: Two Functions**
   - EncryptRanges
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PrefixMatch

- **Property**
  - Answer if a value V matches a range R_i from $[R_1, R_2, ...]$
- **Security**
  - Do not reveal the value of V and R_i
  - If both V_1 and V_2 match R_i, do not reveal the ordering between V_1 and V_2

PrefixMatch vs. OPE

- **Order-preserving Encryption**
  - Preserve the ordering of values after encryption
- **PrefixMatch is better than OPE in this scenario**
  - More secure (No relative ordering)
  - Faster (10000x)
    - Compare with the state-of-the-art OPE schemes (BCLO and mOPE)

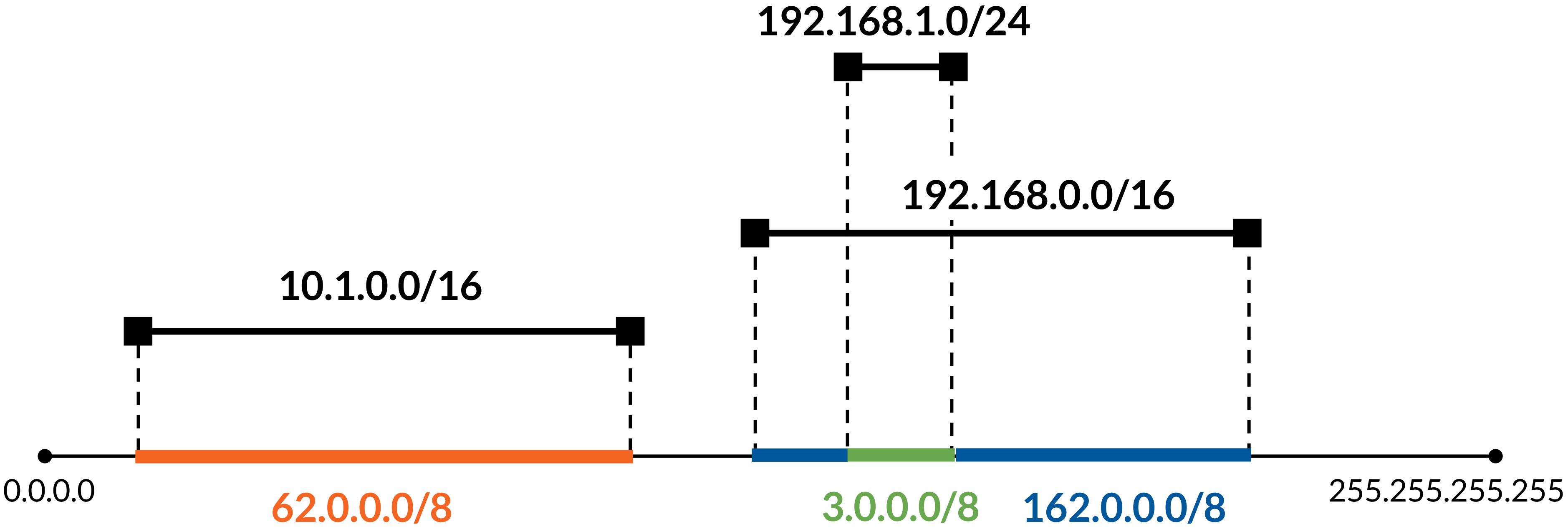
Operation	BCLO	mOPE	PrefixMatch
Encrypt, 10K rules	9333 us	6640 us	0.53 us
Encrypt, 100K rules	9333 us	8300 us	0.77 us
Decrypt	169 us	0.128 us	0.128 us

EncryptRanges

- Firewall Rules

```
block from 192.168.1.0/24 to 205.203.224.0/19
block from 192.168.0.0/16 to 223.254.0.0/16
block from 10.1.0.0/16 to 223.201.0.0/16
```

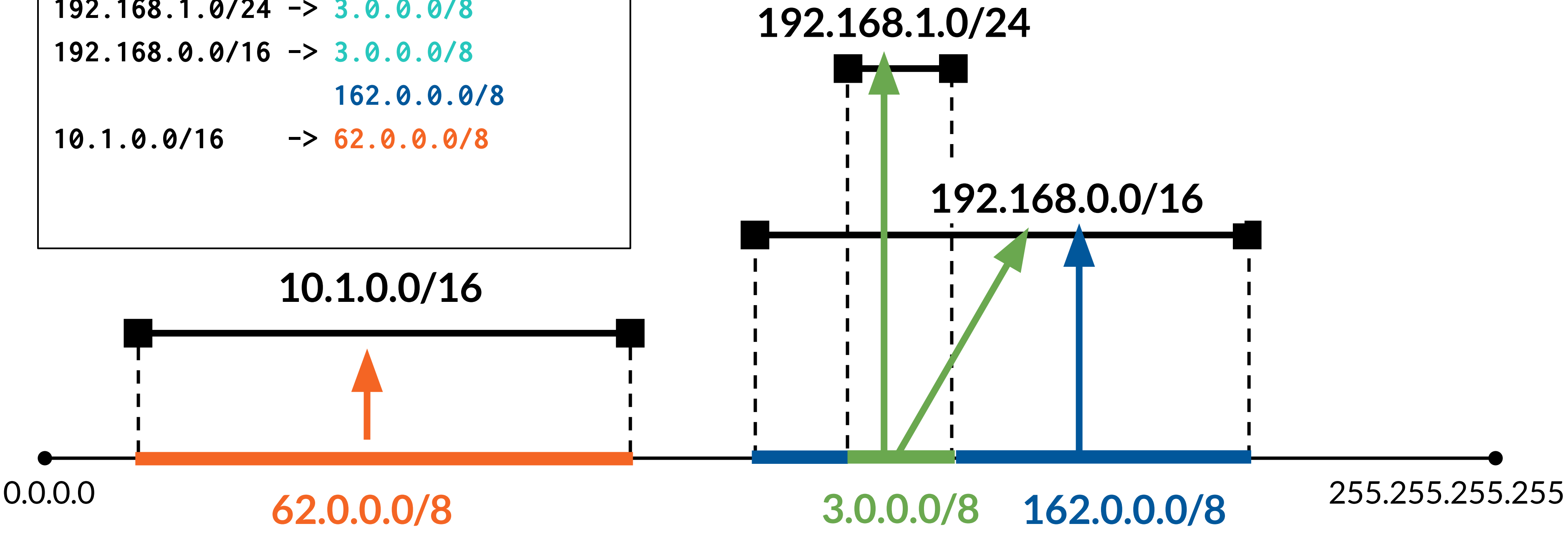
EncryptRanges



Assign Random Prefixes

EncryptRanges

```
192.168.1.0/24 -> 3.0.0.0/8
192.168.0.0/16 -> 3.0.0.0/8
                  162.0.0.0/8
10.1.0.0/16    -> 62.0.0.0/8
```



EncryptRanges

```
block from 192.168.1.0/24 to 205.203.224.0/19
block from 192.168.0.0/16 to 223.254.0.0/16
block from 10.1.0.0/16 to 223.201.0.0/16
```

Source IP

```
192.168.1.0/24 -> 3.0.0.0/8
192.168.0.0/16 -> 3.0.0.0/8
                  162.0.0.0/8
10.1.0.0/16 -> 62.0.0.0/8
```

Destination IP

```
205.203.224.0/19 -> 12.0.0.0/8
223.254.0.0/16 -> 241.0.0.0/8
223.201.0.0/16 -> 163.0.0.0/8
```

```
block from 3.0.0.0/8 to 12.0.0.0/8
block from 3.0.0.0/8 to 241.0.0.0/8
block from 162.0.0.0/8 to 241.0.0.0/8
block from 62.0.0.0/8 to 163.0.0.0/8
```

EncryptValue

- Encrypt each field independently
    - Source IP, Destination IP,

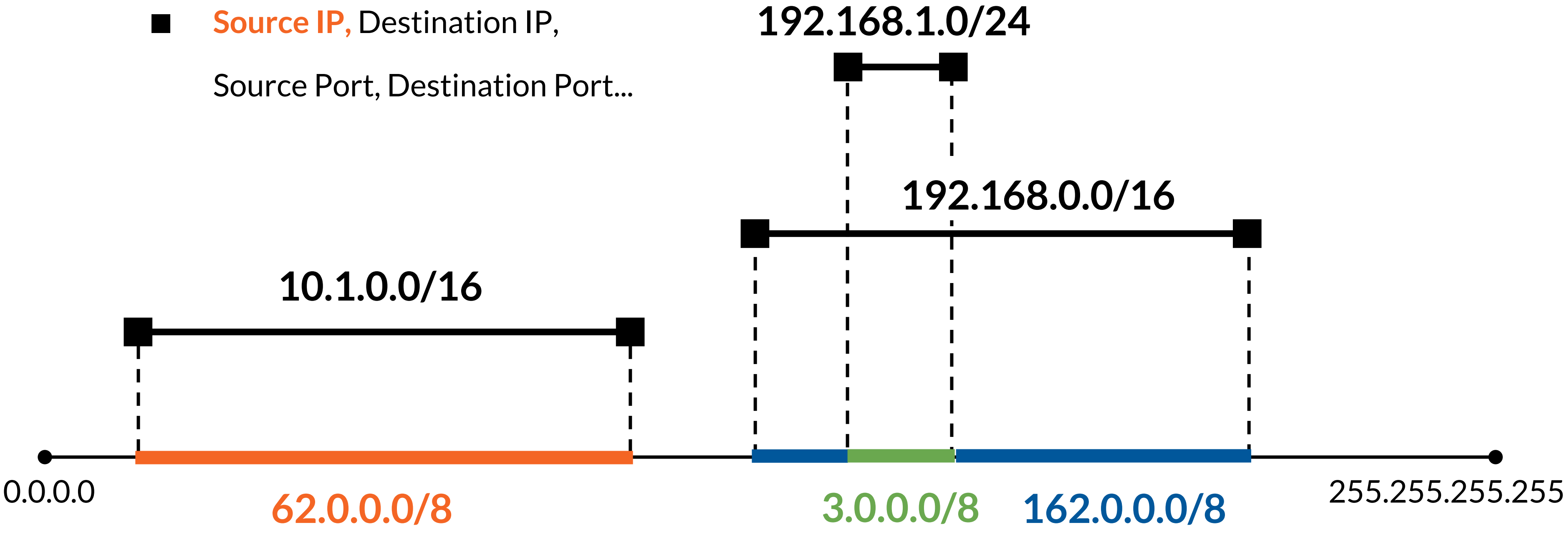
      Source Port, Destination Port...

EncryptValue

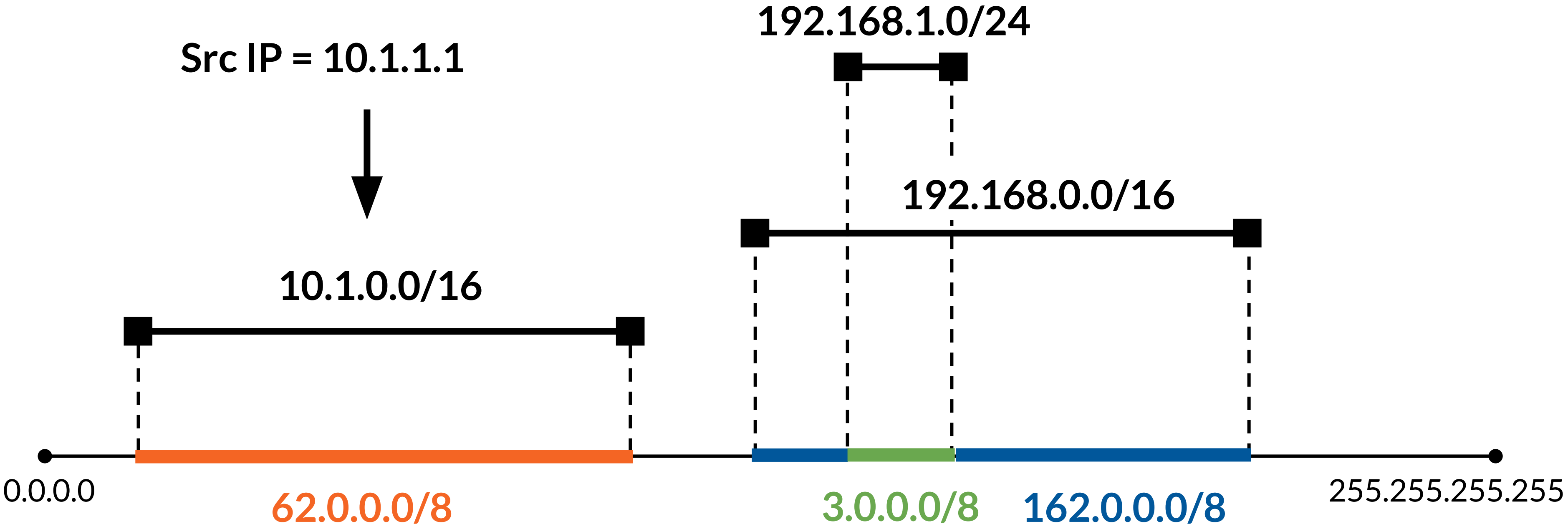
➢ Encrypt each field independently

- Source IP, Destination IP,

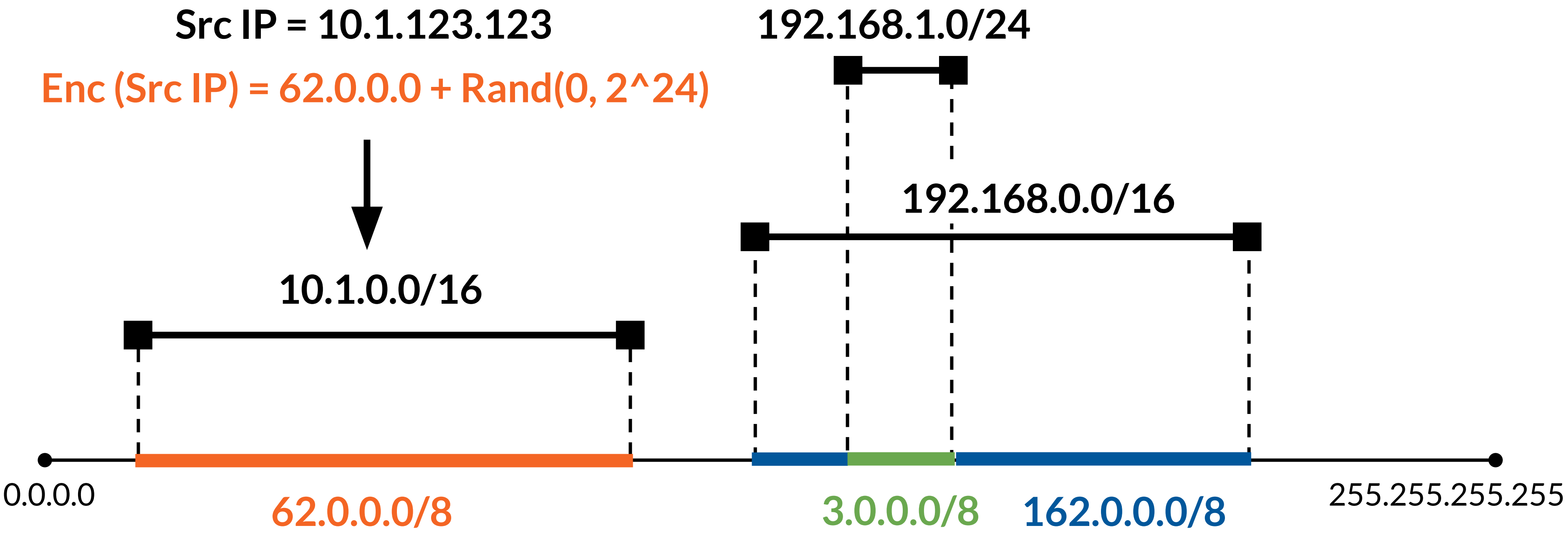
  Source Port, Destination Port...



EncryptValue



EncryptValue



EncryptValue

- **Problem 1: How to support NAT and Load Balancers?**
  - **Deterministic**: The value from the same flow will be mapped to the same value
  - **Injective**: Values from different flows will be mapped to different values
  - **Sufficient condition**

Sufficient condition:

Let

v = (sip, dip, sp, dp, proto)

v' = (sip', dip', sp', dp', proto')

$v = v'$ if and only if

$Enc(v) = Enc(v')$

Src IP = 10.1.123.123

Enc (Src IP) = 62.0.0.0 + Rand(0, 2^24)

EncryptValue

- **Problem 1: How to support NAT and Load Balancers?**
  - Use pseudorandom function,

    seeded by 5-tuple
  - Use IPv6 to avoid collisions

~~Src IP = 10.1.123.123~~
~~Enc (Src IP) = 62.0.0.0 + Rand(0, 2^24)~~

Src IP = ::FFFF:10.1.123.123
Enc (Src IP) = 3e00::/8 + PRF(Src IP)

EncryptValue

- **Problem 1: How to support NAT and Load Balancers?**
- **Problem 2: How to decrypt?**
    - Store AES(Src IP) in IP Options
    - Decrypt AES(Src IP)

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1. Service Model of Embark
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3. **Evaluation**
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Evaluation

- **What kinds of middleboxes does Embark support?**
  - Performance of each type of middleboxes
- **How much does PrefixMatch increase the number of rules?**
- **Microbenchmarks**
  - How does PrefixMatch compare with OPE?
  - How well does PrefixMatch scale with the number of rules?
- **Performance**
  - How fast is the gateway (with PrefixMatch and with KeywordMatch)
  - How much does the service model increase the page load time?

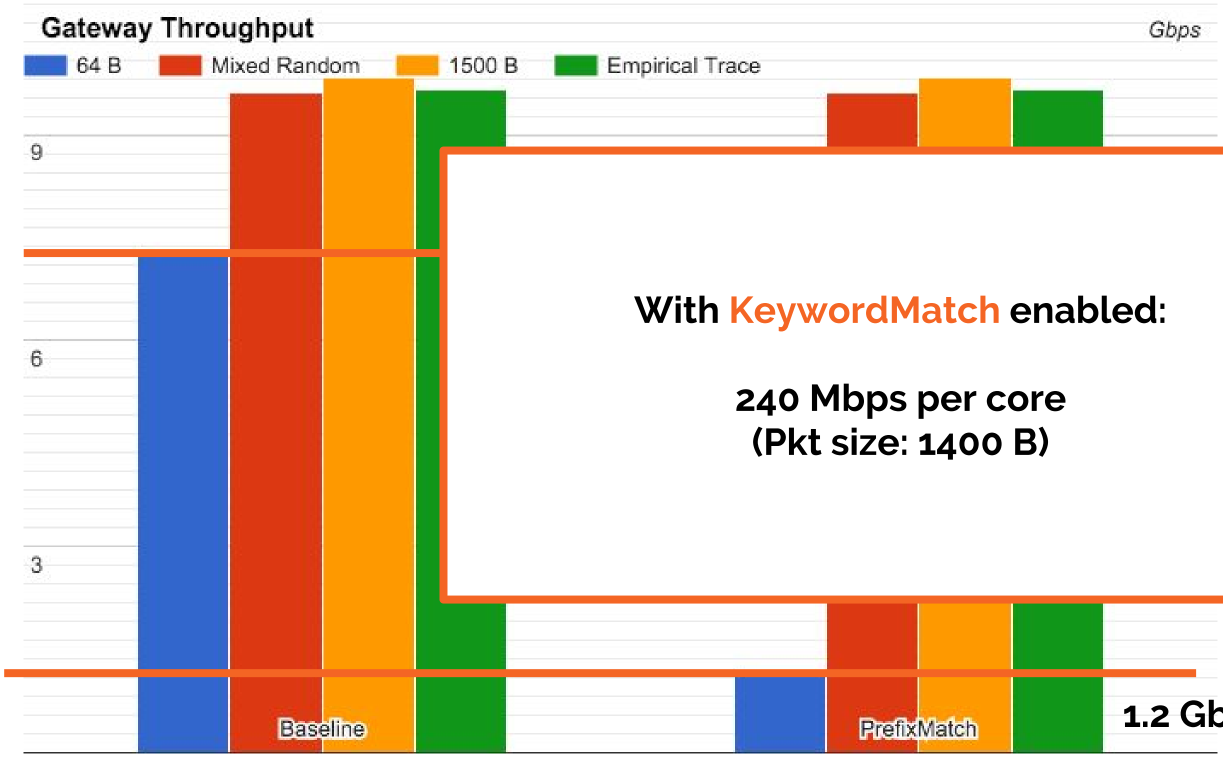
Supported Middleboxes

IP Firewall	Linux iptables	PrefixMatch
NAT	Linux iptables	
L3 Load Balancer	ECMP	
L4 Load Balancer	HAProxy	
HTTP Proxy	Embark vs Squid	KeywordMatch
Parental Filter	Embark vs Squid	
Intrusion Detection (excluding scripts and other statistical techniques)	Embark vs Snort	

How much does PrefixMatch increase Firewall rules?

- **Upper bound**
  - $O(n^d)$, d is the number of fields
- **Empirically**
  - Rulesets
    - 3 firewall rulesets from campus network at UC Berkeley
    - 1 firewall ruleset from Emerging Threats
  - Result
    - UCB rulesets: No increase
    - Emerging Threats: from 1363 to 1370
  - Intuition
    - Most firewall rules don’t overlap

How fast is the gateway (without KeywordMatch)?



1.2 Gbps

Baseline PrefixMatch

With KeywordMatch enabled:

240 Mbps per core
(Pkt size: 1400 B)

rformance with 1k rules:

- **1.2 Gbps** (min-size)
- **Line rate** (other cases)

See the paper for ...

- How we design and implement middleboxes
- Formal proof of sufficient conditions for NAT and L3/TCP Load Balancers
- **Limitations**
- More in-depth evaluation
- ...

Conclusion

Middleboxes can be outsourced in a way that still keeps the traffic confidential with **Embark**.

Paper: changlan.org/papers/embark.pdf

Contact: clan@eecs.berkeley.edu

Thanks!