

Security of Hedged Fiat–Shamir Signatures under Fault Attacks

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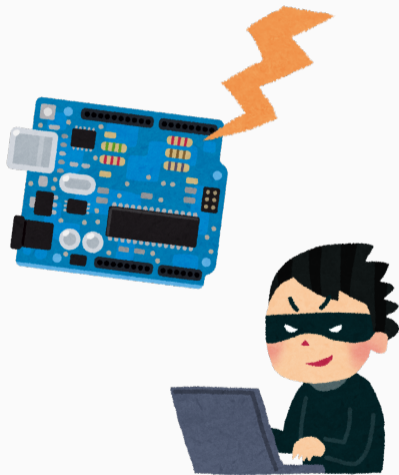
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Goal of Our Work

- Formally analyze the fault-resilience of existing Fiat–Shamir signatures
 - Provable security methodology.
 - Motivated by actual fault attacks on concrete schemes.



1. Randomized signature : $r \leftarrow \text{RNG}(\cdot)$

- Nonces don't need to be uniform: low-quality RNG or counter should suffice.
- Randomness r doesn't repeat on the same message.

To what extent are hedged FS signatures secure against fault attacks?

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3. Hedged signature : $r \leftarrow H(sk, m, nonce)$ 😊 Seems secure?

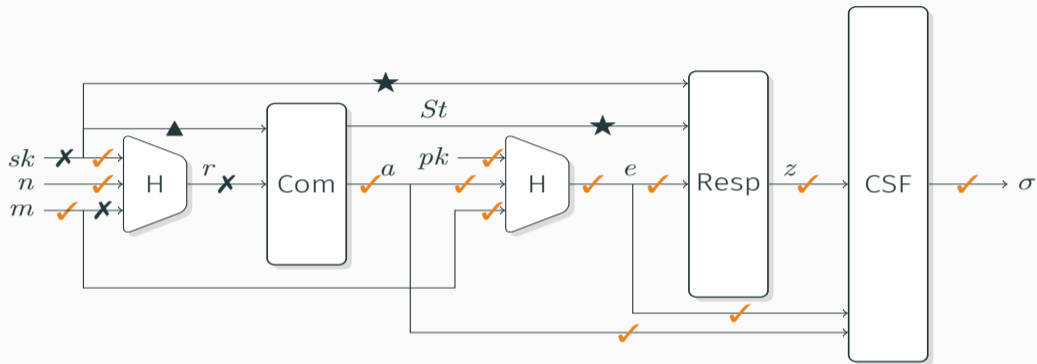
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Contributions

- Formal attacker model and security notions to capture the corrupted nonces and previous fault attacks.
- Proved that **hedged FS schemes in general** are secure against single-bit fault attacks on many intermediate wire values in the signing algorithm.
 - + Negative results for a few wires.
- Application to concrete instantiations.
 - XEdDSA: Hedged variant of EdDSA used in Signal
 - Picnic2: NIST PQC competition round 2 candidate

Overview of Our Results



If \mathcal{A} doesn't query the same (m, n) pair more than once

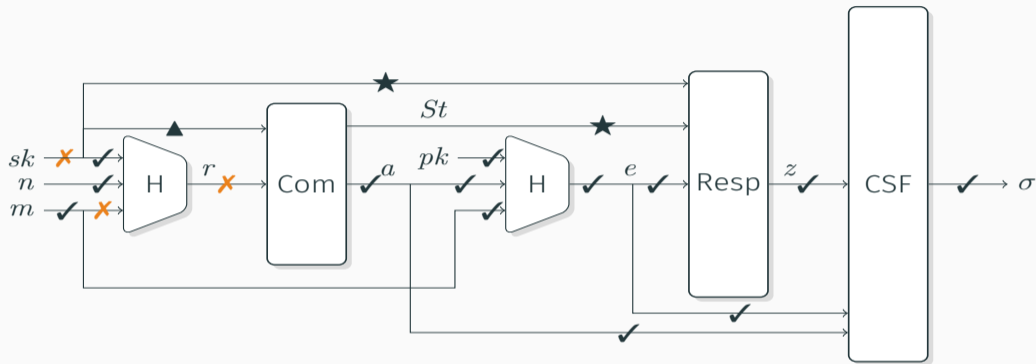
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✗ insecure against single-bit flip/stuck-at faults.

★ security only holds for signatures from subset-revealing ID (e.g., Picnic).

▲ security only holds for signatures from input-delayed ID (e.g., XEdDSA).

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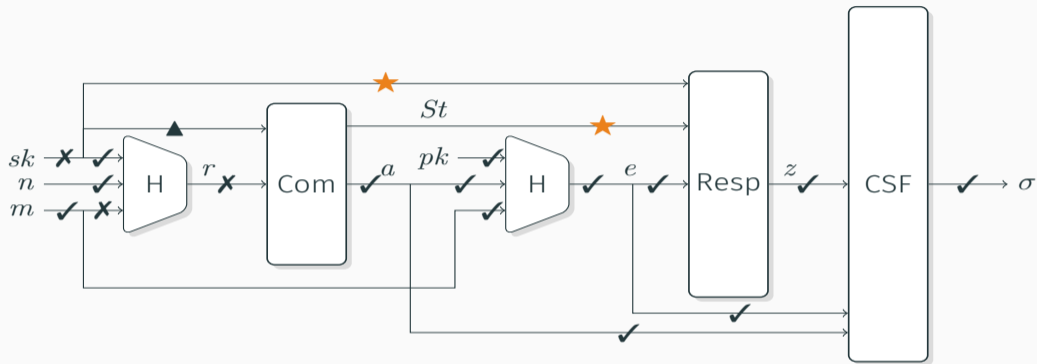
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Conclusion

- Hedged FS is provably more resilient than the randomized/deterministic FS!
 - Negative results show where practitioners pay the most attention.
- Open questions
 - Extension to more advanced fault attacker model.
 - Multi-bit/position faults. Partially handled by Fischlin and Günther (CT-RSA'20) for generic signatures.
 - Fault within Com, Resp or public parameters.
 - Model for instruction skipping faults.
 - Fault + QRROM.
 - Lattice signatures from FS with aborts.

Thank you!

More details at <https://ia.cr/2019/956>